

# MATCHING-PERFECT AND COVER-PERFECT GRAPHS

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## ABSTRACT

It is shown that a graph  $G$  has all matchings of equal size if and only if for every matching set  $\lambda$  in  $G$ ,  $G \setminus V(\lambda)$  does not contain a maximal open path of odd length greater than one, which is not contained in a cycle. ( $V(\lambda)$  denotes the set of vertices incident with some edge of  $\lambda$ .) Subsequently edge-coverings of graphs are discussed. A characterization is supplied for graphs all whose minimal covers have equal size.

## 1. Introduction and definitions

Let  $G = (V, E)$  be a finite graph. A subset of  $E$  is a *matching set* if no two edges of the set are adjacent. A matching set is a *matching* in  $G$  if it is maximal in  $G$  with respect to inclusion. In a finite graph different matchings may have different cardinalities. In [1] Grünbaum uses  $\underline{m}(G)$  and  $\bar{m}(G)$  to denote the smallest and the greatest cardinalities respectively of matchings in  $G$ . In the above-mentioned paper Grünbaum displays all connected graphs  $G$  for which  $\underline{m}(G) = \bar{m}(G) = 2$  and at the same time raises the problem of characterizing those graphs  $G$  for which  $\underline{m}(G) = \bar{m}(G)$ . Let us call such graphs *matching-perfect*, in brief MP.

For  $\lambda \subseteq E$ , let  $V(\lambda)$  denote the set of vertices incident with some edge of  $\lambda$ . If  $V(\lambda) = V$ , then  $\lambda$  is called an *edge-cover* or, in short, a *cover* of  $G$ . A *minimum cover* of  $G$  is a cover of  $G$  having a minimum number of edges. It then follows directly from a result in [2] that a graph  $G$  is MP if and only if every matching of  $G$  is contained in a minimum cover of  $G \setminus I(G)$ , where  $I(G)$  is the set of isolated vertices of  $G$ .

We now suggest a different characterization of MP graphs. We first introduce some definitions. A maximal open path in  $G$  of the form  $(a_0, a_1, \dots, a_n)$  such

that  $(a_0, a_n) \notin E$  is called a *snake* in  $G$ . The *length* of the snake is the number of its edges. A snake is *odd* or *even* according to the parity of its length. If  $\lambda$  is a matching set and  $G \setminus V(\lambda)$  contains a snake  $s$ , then  $\lambda$  is a *blocking* of  $s$  in  $G$ . If  $s$  is a snake in  $G$ , then  $\lambda$  may be taken as  $\emptyset$ .  $s$  itself is referred to as a *blocked* snake.  $|S|$  will denote the number of elements of the set  $S$ .

## 2. Matching-perfect graphs

**THEOREM 1.** *A graph is matching-perfect if and only if it does not contain an odd blocked snake.*

**PROOF.** We show that a graph is not MP if and only if it contains an odd blocked snake. The proof uses the familiar technique of "alternating paths".

Let  $s = (a_0, a_1, \dots, a_k)$  be an odd blocked snake of  $G$ , with  $\lambda$  a blocking of  $s$  in  $G$ . Put

$$\alpha = \{(a_0, a_1), (a_2, a_3), \dots, (a_{k-1}, a_k)\}$$

and

$$\beta = \{(a_1, a_2), (a_3, a_4), \dots, (a_{k-2}, a_{k-1})\}.$$

We see that  $\alpha \cup \lambda$  is a matching set in  $G$  which may be completed by a set  $\mu$  of edges to a matching  $A = \alpha \cup \lambda \cup \mu$  of  $G$ . But then, since neither  $a_0$  nor  $a_k$  are adjacent to any vertex in  $G \setminus V(\alpha \cup \lambda)$ , it follows that  $B = \beta \cup \lambda \cup \mu$  is also a matching of  $G$ . But  $|A| = |B| + 1$ , so  $G$  is not MP.

Conversely, let  $G$  be such that  $\underline{m}(G) < \bar{m}(G)$ , and let  $A$  and  $B$  be matchings of  $G$  with  $|A| > |B|$ . Let  $C$  denote the family of maximal alternating chains in  $G$ , with edges alternatingly in  $A$  and in  $B$ . If all elements of  $C$  were to contain an even number of edges, we would have  $|A| = |B|$ . Hence there is an alternating path  $s = (a_0, a_1, \dots, a_k)$  of odd length  $k$ , that starts with an edge in  $A$ . If  $(a_0, x)$  is an edge of  $G$  with  $x \notin V(\alpha)$ , where  $\alpha$  is defined as above, then  $(a_0, x) \notin A \cup B$  since  $A$  is a matching and  $s$  is maximal, but  $x \in V(B)$  since  $B$  is a matching; similarly for  $a_k$ . Let  $B^*$  denote the edges of  $B$  not in  $s$ ; then it is clear that  $s$  is an odd blocked snake in  $G$ , with  $B^*$  a blocking of  $s$  in  $G$ .

The characterization obtained above, although not descriptive, suggests a method to detect a non-matching-perfect graph, which in many cases proves quite efficient.

## 3. Cover-perfect graphs

Let  $G$  be a graph without isolated vertices. Let  $\underline{M}(G)$  and  $\bar{M}(G)$  be the smallest and the greatest number of edges, respectively, in a minimal cover of

$G$ . Let those graphs  $G$  for which  $\underline{M}(G) = \bar{M}(G)$  be called *cover-perfect*, CP in short. We shall assume henceforth that  $G$  has no isolated vertices. A subset  $\lambda$  of  $E$  is *essential* in  $G$  if  $G \setminus \lambda$  is without isolated vertices and  $\bar{m}(G \setminus \lambda) < \bar{m}(G)$ . We then have the following characterization of CP graphs.

**THEOREM 2.** *A graph  $G$  is cover-perfect if and only if it has no essential set of edges.*

This theorem follows quite easily from [2, Th. 2(i)].

It is quite easy to see that an MP graph has no essential edge. It does not follow, however, that an MP graph is necessarily CP. For example the complete graph  $K_n$  is MP for all  $n$ , but is CP only if  $n < 4$ . However, when  $G$  is a tree, we have:

**THEOREM 3.** *If a tree is matching-perfect, then it is cover-perfect.*

**PROOF.** Let  $L(v)$  denote the set of edges incident with the vertex  $v$ . We first prove a lemma.

**LEMMA.** *If  $G$  is a tree and  $\lambda$  a subset of  $E$  such that for every vertex  $v$  of  $G$  we have  $L(v) \not\subseteq \lambda$ , then there is a matching set  $\eta$  in  $G$  such that  $\eta \cap \lambda = \emptyset$  and  $V(e) \cap V(\eta) \neq \emptyset$  for every  $e \in \lambda$ . Moreover there is an  $\eta$  for which  $|\eta| = |\lambda|$ .*

**PROOF.** Let  $x_0$  be an end vertex of  $G$ . For each  $(y, y') \in \lambda$ , if  $y'$  disconnects  $G$  between  $x_0$  and  $y$ , we choose an  $x$  such that  $(x, y) \notin \lambda$ , and include  $(x, y)$  in  $\eta$ . Clearly  $\eta$  thus constructed satisfies the requirements of the lemma.

Now suppose the tree  $G$  is MP and not CP. Then by Theorem 2,  $G$  has an essential set. Denote it by  $\lambda$ .  $\lambda$  satisfies the conditions of the lemma and hence there is a matching set  $\eta$  in  $G$  such that  $\eta \cap \lambda = \emptyset$  and  $V(e) \cap V(\eta) \neq \emptyset$  for every  $e \in \lambda$ . Complete  $\eta$  to a matching  $\eta_0$  in  $G$ . Since  $G$  is MP, we have  $|\eta_0| = \bar{m}(G) \leq \bar{m}(G \setminus \lambda)$ . Then  $\lambda$  is not essential in  $G$ , contrary to our assumption. This proves the theorem.

**COROLLARY.** *Let  $G$  be a tree and  $\lambda$  an essential set in  $G$ . Then  $|\lambda| < \frac{1}{2}|E|$ .*

## REFERENCES

1. B. Grünbaum, *Matchings in polytopal graphs*, Networks (to appear).
2. M. Lewin, *A note on line coverings of graphs*, Discrete Math. **5** (1973), 283–285.

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